Synthesizing Protocol Specifications from Service Specifications in Timed Extended Finite State Machines  *

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Abstract

We propose a specification model and present a method to algorithmically derive a protocol specification from a service specification based on the model. Unlike the previous models based on finite state machines, the proposed model can explicitly express concurrency, synchronization, and timing requirements such as delays and timeouts. We assume that there exists a reliable communication channel between any two protocol entities and the maximum delay for each channel is bounded by a positive constant. Because of the variable nature of the communication delays along with the time constraints associated with events, no protocol specification can fully simulate the service specification. The proposed method derives a protocol specification that is optimal in the sense that it provides the largest possible subset of the service specification under the communication delay constraints. We also give a method to derive a sub specification from a service specification and a maximum communication delay of each channel such that the sub specification, but no superset of it, can be simulated by the derived protocol specification.

1 Introduction

There are two common approaches for designing communication protocols: analysis and synthesis [4]. In the analysis method, the protocol designer begins with a preliminary version of the protocol usually obtained by ad hoc methods. This approach usually results in an incomplete and erroneous design, which is followed by an analysis and redesign process. The sequence of redesign, analysis, and error correction is applied iteratively until an error-free design is obtained. In the synthesis method, a partially specified or incomplete protocol design is completed incrementally, or automatically, without any interaction by the designer such that as the synthesis process proceeds correctness is maintained. The process ends with a design that provides the set of specified services. Therefore, no further verification of the protocol design is necessary as in the analysis approach.

Much research has been done in the area of protocol synthesis. The reader may refer to [1] for a survey and assessment of several synthesis methods. The synthesis methods can be

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classified by the modeling formalism. The models include finite state machines [2, 3, 7], Petri nets [5], LOTOS-like models [8, 9], etc. However, these methods do not provide or represent the notion of time, which is important for the proper functioning of communication systems.

Recently, a few methods [10, 11] have been proposed that derive protocol specifications from timed service specifications. In [10], a model based on finite state machines has been proposed for specifying timing requirements by using a global clock, timers, and counters. The method derives the protocol and medium specifications from a service specification written as a set of timed transitions. The model represents temporal requirements between remote as well as consecutive events, which necessarily introduces an exponential increase in the number of timers. On the other hand, [11] has proposed a model based on LOTOS that restricts the time constraints of service specifications while fixing the maximum delay of the communication media in the sense that the model can specify a complicated order of events in a structural way.

In this paper, we propose a model called timed extended finite state machine (TEFSM) based on the extended finite state machine (EFSM) model to deal with timed operations between consecutive events. Delay and timeout are certainly two of the most useful timed operations. To represent these events, we use the notion of a timed transition in our model by associating a time interval $[l, u]$ with the transition. The lower and upper time limits are measured with respect to a global clock, and can thus be used in modeling timed properties including delays and timeouts. The notion of a timed transition is not new, and our model is in fact inspired by a few previous works [12, 13]. The main difference is that our model can express concurrency and synchronization among protocol entities explicitly while these previous models could not. For synthesis, we assume that each communication channel is error-free and has a propagation delay bounded by a constant, as in [11]. We present an algorithm that derives a protocol specification from a service specification modeled as a TEFSM when an upper bound of delay for each channel is given.

The paper is organized as follows. Section 2 describes our TEFSM model. Section 3 formalizes the protocol synthesis problem and gives some notation. In section 4, we present an algorithm for deriving the protocol specification of a protocol entity from a service specification and prove the correctness of the algorithm by investigating the relationship between the service specification and the protocol specification. In section 5, we demonstrate the applicability of our synthesis method by giving an example. Section 6 gives some concluding remarks and discusses areas requiring future work.

## 2 The Model

The TEFSM model is designed as a method for the formal description of service and protocol specifications. A TEFSM $M$ is defined by a tuple $< S, FJ, V, T, s_0 >$, where (1) $S$ is a nonempty set of states and for each $s \in S$, $s$ is a choice, fork, join, or fork/join state. To represent a possible parallel execution among protocol entities, $M$ explicitly uses a pair $(fork, join)$ of states such that the control flow (directed paths) from fork to join can be executed concurrently and independently. If a join state is also a fork state which is matched with another join state, we call it fork/join state. Note that for each fork state, there exists a unique join state and vice versa. All states other than fork, join, or fork/join states are choice states. If more than one outgoing transition exists for a choice or a join state, $M$ can arbitrarily choose one transition and execute it. See Figure 1 for an example of the classification of states; (2) $FJ$ is
a finite (possibly empty) set of (fork, join) pairs in \( M \); (3) \( V \) is a set of variables including input, output, and local variables, denoted by \( I, O, \) and \( L \), respectively; (4) \( T \) is a set of transitions and each transition \( t \in T \) is a 6-tuple \( \langle \text{head}(t), \text{tail}(t), P(t), E(t), \text{host}(t), [\text{min}_t, \text{max}_t] \rangle \), where \( \text{head}(t) \) and \( \text{tail}(t) \) are respectively the head and the tail state of \( t \), \( P(t) \) is the enabling predicate in \( V \) associated with \( t \), \( E(t) \) is the event on \( V \) associated with \( t \), \( \text{host}(t) \) is the protocol entity that executes \( t \), and \( [\text{min}_t, \text{max}_t] \) is the time interval associated with \( t \) such that \( t \) can be executed after \( \text{min}_t \), but no later than \( \text{max}_t \) has passed since the time of visit to \( \text{head}(t) \) state. If a time interval is not specified explicitly, the default interval \([0, \infty)\) is assumed. An event is a partial function: \( E(t) : L \times I \rightarrow L \times O \). We denote \( a^t[\text{min}_t, \text{max}_t] \) as the transition \( t \) with the action \( a \) and the protocol entity \( i \) when the other components of \( t \) are of no concern; (5) \( \delta \) is a partial state transition function such that \( \delta : S \times T \rightarrow S \), and (6) \( s_0 \) is an initial state.

![Diagram of a TEFSM](image)

**Figure 1: The Classification of States: An Example**

The execution of a transition \( t \) is an instantaneous action in which both the event associated with \( t \) and the state change to the tail state of \( t \) occur simultaneously. A transition \( t \) in a TEFSM \( M \) must be executed within its time interval if (1) \( M \) is in \( \text{head}(t) \); (2) A finite time interval is associated with \( t \); and (3) \( t \) is enabled throughout the time interval.

A protocol is specified as a set of processes \( \langle PS_0, PS_1, ..., PS_n \rangle \) where each process \( PS_i \) is a TEFSM that can communicate with other processes through FIFO channels. Note that each process \( PS_i \) has only choice states since no concurrent execution is allowed.

A channel from \( PS_i \) to \( PS_j \) has a maximum delay \( D_{ij} \) such that the message transmissions are carried out within \( D_{ij} \), i.e., \( 0 < \text{delay}(i, j) \leq D_{ij} \).

A sending transition \( s_{ij}(m) \) denotes the nonblocking transmission of the message \( m \) from \( PS_i \) to \( PS_j \) and a receiving transition \( r_{ij}(m) \) denotes the blocking reception of the message \( m \) coming from \( PS_i \) to \( PS_j \). Note that \( s_{ij}(m) \) and \( r_{ij}(m) \) are dual events.

Since a TEFSM can be described as a labeled directed graph, we will use state and node, and event and transition interchangeably for the rest of the paper.

### 3 Synthesis Problem

We assume that there is a global digital clock that ticks at a constant frequency and all of the relative times of the protocol entities refer to this clock.

**Notation 1**

(1) Given a finite sequence \( \sigma \), \( \text{first}(\sigma) \) and \( \text{last}(\sigma) \) denote the first and the last element of \( \sigma \), respectively. Denote \( \cdot \) for concatenation. \( \varepsilon \) denotes an empty sequence, \( |\varepsilon| = 0 \).

(2) Given a sequence of events \( \sigma \), we denote \( \sigma_{\leq i} \) for the projection of \( \sigma \) onto the events of \( PS_i \).
(3) For a state \( s \) in a TEFSM, \( t(s) \) denotes the time when the machine has visited the state.
(4) Given a state \( s \) in a TEFSM, \( IN(s) \) and \( OUT(s) \) denote the sets of the incoming and the outgoing transitions of \( s \), respectively.

**Definition 1** For a join or a fork/join state \( s \) in a TEFSM, \( t(s) \) is defined as \( \max_{1 \leq i \leq k} \{ t_i \} \) where \( t_i \) is the time when the incoming event \( e_i \) of \( s \) has occurred, \( 1 \leq i \leq k \). For all other states \( s \) in a TEFSM, \( t(s) \) is defined to be the time when an incoming event \( e \) of \( s \) has occurred provided the machine has executed \( e \) to reach the state \( s \).

**Definition 2** [10] A timed sequence \( S \) in a TEFSM \( M \) is a finite or infinite sequence of pairs \( < e_i, t_i > \), where \( t_i < t_{i+1} \) if \( host(e_i) = host(e_{i+1}) \) and \( t_i \leq t_{i+1} \) otherwise and each pair \( < e_i, t_i > \) denotes that an event \( e_i \) of \( M \) has occurred when the time is equal to \( t_i \).

**Definition 3** A timed sequence \( S \) in a TEFSM \( M \) is valid if each \( e_i \) has been executable at \( head(e_i) \), i.e., \( P(e_i) \) was true at \( head(e_i) \) and has remained to be true till the execution of \( e_i \), and \( t(head(e_i)) + \min_{i} t_i \leq t(\text{head}(e_i)) + \max_{i} t_i \).

**Definition 4** Let \( \{seq_1, \ldots, seq_n\} \) be a set of sequences, where \( seq_i \) is a valid sequence in a TEFSM \( PS_i, 1 \leq i \leq n \). A merged sequence \( seq(\sum_{i=1}^{n} i) \) from the set is a sequence of pairs \( < e_i, t_i > \), where \( e_i \) is in the union of the events of \( PS_j, 1 \leq j \leq n \), such that \( seq(\sum_{i=1}^{n} i) \mid j = seq(j) \), for each \( j, 1 \leq j \leq n \), and \( t_i \leq t_{i+1} \).

**Notation 2** (1) \( \{PS_i, 1 \leq i \leq n\} \) denotes the set of the merged sequences \( \{seq(\sum_{i=1}^{n} i)\} \) is a merged sequence from \( \{seq_1, \ldots, seq_n\} \), where \( seq_i \) is a valid sequence in a TEFSM \( PS_i, 1 \leq i \leq n \). (2) \( \{SS\} \) denotes the set of valid sequences in a service specification \( SS \).

The protocol synthesis problem is basically to derive a protocol specification for the protocol entities from a given service specification such that each protocol entity would be able to execute events in exactly the same order as specified in the service specification. However, since the specification is modeled by a TEFSM, the problem now is to consider time constraints as well as the relative order of the events in the service specification. Along with the time constraints associated with events, the variable nature of the communication delays make it impossible to derive a protocol specification which would be able to fully simulate the service specification. Therefore, to cope with the discrepancy between protocol and service specifications, we define the protocol synthesis problem as follows. Derive a protocol specification from a given service specification which satisfies the following conditions.

**Definition 5** A derived protocol specification \( PS_i, 1 \leq i \leq n \), is correct with respect to the service specification \( SS \) if (1) every merged sequence \( seq(\sum_{i=1}^{n} i) \) from \( \{seq(1), \ldots, seq(n)\} \), where \( seq(i) \) is a valid sequence in \( PS_i, 1 \leq i \leq n \), is a valid sequence in \( \{SS\} \); and (2) every valid sequence \( \sigma \) in \( \{SS\} \) is a merged sequence from \( \{SS \mid 1, \ldots, SS \mid n\} \), where \( SS \mid i \) preserves the order of events as specified in \( PS_i, 1 \leq i \leq n \).

Condition (2) of Definition 5 means that the derived protocol specification should preserve the order of events, but not necessarily simulate the same time stamp of the events in the service specification.
4 Synthesis Algorithm

We present an algorithm that derives the maximal protocol specification among the correct protocol specifications from a service specification. Moreover, we also give an algorithm for finding the maximal subset of a service specification which can be represented by the derived protocol specification.

Since we assume that each protocol entity is modeled by a TEFSM, no (fork, join) pair in a service specification should contain a set of control flows that might be able to cause a conflict, i.e., two or more concurrent events with the same host (protocol entity) and the same time stamp. To cope with the problem, we provide a sufficient condition for a service specification to be conflict-free. We believe that the condition given in Lemma 1 does not severely restrict the modeling power of TEFSM.

Lemma 1 A TEFSM $M$ with nonempty $FJ$ is conflict-free if for each $(f, j) \in FJ$, any two sequences $s_1$ and $s_2$ from $f$ to $j$ that can be executed concurrently by $M$ do not share a host, i.e., $host(s_1) \cap host(s_2) = \emptyset$, where $host(s_i) = \{m|a^m\text{ is an event in }s_i\}$.

We also impose a restriction $R_1$ to the service specification $SS$ as follows: for every choice state $s$ in $SS$, $|host(OUT(s))| = 1$. $R_1$ means that when a choice is possible during the execution of a concurrent protocol system, the choice should be made locally by the same protocol entity to avoid possible deadlocks.

For the sake of algorithm presentation, we denote $(x, e^i, y)$ as the event $e^i$ with the head state $x$ and the tail state $y$, respectively. The following algorithm generates $PS_i$, the specification for protocol entity $i$, and we can get the protocol specification $PS_i, 1 \leq i \leq n$, by running the algorithm $n$ times with different $i$ each time.

Synthesis

- Input: Service specification $SS$ with the condition in Lemma 1 and $R_1$ represented by a TEFSM and $D_{i,j}, \forall i, j, 1 \leq i, j \leq n$. Note that $D_{i,i} = 0, \forall i, 1 \leq i \leq n$.
- Output: Protocol entity specification $PS_i$ in a TEFSM

1. For each state $s$ with $|IN(s)| > 0$ in $SS$ do the following:
   Let $IN(s) = \{(u_1, e^{in_1}, s), \ldots, (u_k, e^{in_k}, s)\}$ and $OUT(s) = \{(s, f^{out_1}, v_1), \ldots, (s, f^{out_l}, v_l)\}$.
   (a) (Append send and/or receive transitions appropriately.)
      i. $s$ is a choice state: Note that $out_1 = \ldots = out_l = j$.
         for each transition $(u_x, e^{in_x}, s), 1 \leq x \leq k,$ do:
         - if $(in_x \neq i \land j = i)$, then append a receive transition to the transition as in Figure 2(a);
         - else if $(in_x = i \land j \neq i)$, then append a send transition to the transition as in Figure 2(b);
         - else if $(in_x = i \land j = i) \lor (in_x \neq i \land j \neq i)$, then do nothing;
      ii. $s$ is a fork, but not a join state:
         for each transition $(u_x, e^{in_x}, s), 1 \leq x \leq k,$ do:
         - if $(in_x = i)$, then append a set of send transitions to the transition as in Figure 3(a);
Figure 2: Case(i) $s$ is a choice state

- else if $(in_x \neq i \land i \in \{out_1, \ldots, out_l\})$, then append a receive transition to the transition as in Figure 3(b);
- else if $(in_x \neq i \land i \notin \{out_1, \ldots, out_l\})$, then do nothing;

![Diagram](a)

Figure 3: Case(ii) $s$ is a fork, but not a join state

iii. $s$ is a join or a fork/join state:
- if $(i \in \{in_1, \ldots, in_k\}) \land (i \in \{out_1, \ldots, out_l\})$, then append a set of send and receive transitions to the transition as in Figure 4(a);
- else if $(i \in \{in_1, \ldots, in_k\}) \land (i \notin \{out_1, \ldots, out_l\})$, then append a set of send transitions to the transition as in Figure 4(b);
- else if $(i \notin \{in_1, \ldots, in_k\}) \land (i \in \{out_1, \ldots, out_l\})$, then append a set of receive transitions to the transition as in Figure 4(c);
- else if $(i \notin \{in_1, \ldots, in_k\}) \land (i \notin \{out_1, \ldots, out_l\})$, then do nothing;

(b) Adjust the time intervals associated with the outgoing transitions of $s$, if necessary.)

if $i \in \{out_1, \ldots, out_l\}$, then [for each transition $t^y = (s, f^{out_y}, v_y)$, $1 \leq y \leq l$, such that $out_y = i$, do: $[\text{min}_{t^y}, \text{max}_{t^y}] \leftarrow [\text{min}_{t^y}, \text{max}_{t^y} - \max_{1 \leq \varepsilon \leq k} \{D_{in_{s,\varepsilon}}\}]$]

2. (Project the $SS$ from the Step 1 onto $PS_i$.) For every event $e^z, z \neq i$, replace the event with an $\epsilon$ transition.

3. For each pair $(s_f, s_j) \in FJ$, remove all $\epsilon$ paths from $s_f$ to $s_j$, if any. If all the transitions and states except $s_f$ and $s_j$ are removed, then merge $s_f$ and $s_j$ into a single state $s_{f,j}$.

4. Remove $\epsilon$ transitions by the standard algorithm given in [6].

\[\text{If } \min_{t^y} > \max_{t^y} - \max_{1 \leq \varepsilon \leq k} \{D_{in_{s,\varepsilon}}\}, \text{ then the execution of } f^{out_y} \text{ may not be possible under the delay constraint.}\]
Lemma 2 Let $PS_i, 1 \leq i \leq n$, be the derived protocol specification from $SS$ under the delay constraints $D_{i,j}, 1 \leq i, j \leq n$. Then $\{PS_i, 1 \leq i \leq n\} \subseteq \{SS\}$. Moreover, $\{PS_i, 1 \leq i \leq n\}$ is maximal in the sense that any extension of a time interval in any $PS_i$ might be able to generate sequences which are not in $\{SS\}$ under some specific delay constraints.

Proof: We first show that $\{PS_i, 1 \leq i \leq n\} \subseteq \{SS\}$. Let $\{seq(i), 1 \leq i \leq n\}$ be a set of sequences such that $seq(i)$ is a valid sequence in $PS_i$ for each $i, 1 \leq i \leq n$. The proof is by induction on $|\sigma|$, where $\sigma$ is a merged sequence from $\{seq(i), 1 \leq i \leq n\}$. Base Case $|\sigma| = 1$. It is clear that $\sigma$ is a valid sequence in $SS$. Induction Hypothesis (IH for short) Assume the claim holds for $|\tau| = k > 0$. Let $\sigma' = \sigma < a^i, t > |\tau| = k$, be a merged sequence from $\{seq(i), 1 \leq i \leq n\}$. Let $s = head(a^i)$. Suppose $s$ is a choice state. By the Step 1(a) of the algorithm Synthesis, $PS_i$ can execute $a^i$ either after having received a message from one of $PS_{in_x}, in_x \neq i$ or after having executed $e^{in_x}, in_x = i$. In either case, we know from the algorithm Synthesis that an event in $IN(s)$ must have occurred in $\sigma$. Let $e^{in_x}$ be the latest event from $IN(s)$ in $\sigma$. Then, the subsequence $e^{in_x}, \ldots, last(\sigma)$ of $\sigma$ does not have any event from $OUT(s)$ since otherwise $a^i$ would not have occurred in $\sigma$ by the nature of choice state. Thus, we showed that $a^i$ had been executable at $last(\sigma)$. Suppose $s$ is a fork, but not a join state. The only difference here from the above case (a choice state) is that the subsequence $e^{in_x}, \ldots, last(\sigma)$ of $\sigma$ might have some events from $OUT(s)$, but no $a^i$'s since otherwise $a^i$ would not have occurred at $last(\sigma)$. Suppose $s$ is a join, but not a fork state. By the construction of $PS_i$ in Step 1(a) of the algorithm Synthesis, we know that $PS_i$ can execute $a^i$ only after having received a set of messages from $\{PS_{in_x}\}$ where $(u_x, e^{in_x}, s) \in IN(s), in_x \neq i$ , which implies that $\{e^{in_x}, in_x \neq i\}$ had occurred in $\sigma$ by $PS_{in_x}$, respectively. Also, we know that $e^{in_y}$, where $(u_y, e^{in_y}, s) \in IN(s)$, and $in_y = i$, if any, had occurred in $\sigma$. Let the most recently occurred event from $OUT(s)$ in $\sigma$ be $e^{in_y}$, i.e., $t(s)$ is equal to the time when $e^{in_y}$ has occurred. Then the subsequence $e^{in_y}, \ldots, last(\sigma)$ of $\sigma$ does not have any event from $OUT(s)$ since otherwise $a^i$ would not have occurred. Thus, $a^i$ had been executable at $last(\sigma)$. Suppose

Figure 4: Case(iii) $s$ is a join or a fork/join state
s is a join/fork state. As above, we know that \( \{e^{in_s}, 1 \leq x \leq k \} \) had occurred in \( \sigma \). We also know that the subsequence \( e^{in_s}, \ldots, \) last(\( \sigma \)) might have some events from \( OUT(s) \), but no \( a^i \)'s since otherwise \( a^i \) would not have occurred at \( \text{last}(\sigma) \), where \( e^{in_s} \) is the most recently occurred event from \( OUT(s) \) in \( \sigma \). Now, it is straightforward that \( a^i \) had been executable at \( \text{last}(\sigma) \). Thus, we conclude that \( a^i \) had been executable at \( \text{last}(\sigma) \) for all cases. Next we show that \( t(s) + \min_{a^i} \leq t \leq t(s) + \max_{a^i} \), where \( [\min_{a^i}, \max_{a^i}] \) is the time interval associated with \( a^i \) in SS. The time interval associated with \( a^j \) in \( PS_i \) by the algorithm Synthesis Step 1(b), becomes \( [\min_{a^j}, \max_{a^j} - \max_{1 \leq x \leq k}\{D_{in,j}\}] \). Since \( \sigma' |_{i} = \sigma |_{i} \cdot < a^j, t > \) is a valid sequence in \( PS_i \), we know that \( t(s) + \min_{a^j} + d_{in,j} \leq t \leq t(s) + \max_{a^j} - \max_{1 \leq x \leq k}\{D_{in,j}\} \), where \( 0 < d_{in,j} \leq D_{in,j} \). Thus, we have that \( t(s) + \min_{a^j} < t(s) + \min_{a^j} + d_{in,j} \leq t \leq t(s) + \max_{a^j} - \max_{1 \leq x \leq k}\{D_{in,j}\} \). Therefore, since \( \sigma \) is a valid sequence in SS by III, \( a^j \) is also a valid sequence in SS from the above argument. To prove the maximality of \( \{PS_i, 1 \leq i \leq n\} \), consider a sequence \( \psi \) in \( \{SS\} - \{PS_i, 1 \leq i \leq n\} \). It is clear that \( |\psi| > 1 \), since any sequence in \( \{SS\} \) with length 1 should also be in \( \{PS_i\} \), for some \( i \). We know that there exists a pair of \( s < e^i, t_e >, < f^i, t_f > \) in \( \psi \) such that \( e^i \in IN(s), f^i \in OUT(s), i \neq j \), and \( t(s) = t_e \) for some state \( s \) in SS, since otherwise \( \psi \) would not be in \( \{SS\} - \{PS_i, 1 \leq i \leq n\} \). Note that \( e^i \) and \( f^j \) might not be adjacent in \( \psi \). By the algorithm Synthesis, the time interval associated with the event \( f^j \) is adjusted into \( [\min_{f^j}, \max_{f^j} - \max_{1 \leq x \leq k}\{D_{in,j}\}] \) in \( PS_j \), where \( e^{in_j} \in IN(s), 1 \leq x \leq k \). Assume the interval associated with the event \( f^j \) in \( PS_j \) is extended to \( [\min_{f^j} - \epsilon_j, \max_{f^j} - \max_{1 \leq x \leq k}\{D_{in,j}\} + \epsilon_j] \), where \( \epsilon_j \) and \( \epsilon_j \) are positive constants. Then \( t_e + \min_{f^j} - \epsilon_j + \eta < t_e + \min_{f^j} \) for a positive constant \( \eta \) such that \( \eta < \epsilon_j \leq D_{i,j} \). Hence, if the actual delay from \( PS_i \) to \( PS_j \) is \( \eta \), then \( < e^i, t_e >, < f^j, t_e + \min_{f^j} - \epsilon_j + \eta > \) would not be possible in any sequence in \( \{SS\} \), since \( t_e + \min_{f^j} - \epsilon_j + \eta \leq t_e + \min_{f^j} \). Similarly, \( t_e + \max_{f^j} - \max_{1 \leq x \leq k}\{D_{in,j}\} + \epsilon_j + \max_{1 \leq x \leq k}\{D_{in,j}\} \) would not be possible in any sequence in \( \{SS\} \), if \( \max_{1 \leq x \leq k}\{D_{in,j}\} = D_{i,j} \) and the actual delay from \( PS_i \) to \( PS_j \) is \( D_{i,j} \). Note that if \( \max_{1 \leq x \leq k}\{D_{in,j}\} > D_{i,j} \), the validity of the pair \( < e^i, t_e >, < f^j, t_e + \max_{f^j} - \max_{1 \leq x \leq k}\{D_{in,j}\} + \epsilon_j + D_{i,j} > \) in \( SS \) depends upon the sign of the value \( D_{i,j} + \epsilon_j - \max_{1 \leq x \leq k}\{D_{in,j}\} \).

On the other hand, it should be clear that \( \{SS\} \not\subseteq \{PS_i, 1 \leq i \leq n\} \) because of the adjustment in Step 1(b) of the algorithm Synthesis. However, we can restrict \( SS \) to get a sub specification \( SS^* \) such that \( \{SS^*\} \subseteq \{PS_i, 1 \leq i \leq n\} \). Here, we give an algorithm to generate such a sub specification \( SS^* \) which is maximal in the sense that \( \{SS^*\} \subseteq \{PS_i, 1 \leq i \leq n\} \) implies \( \{SS^*\} \subseteq \{SS^*\} \).

**Restriction**

- **Input:** Service specification \( SS \) with the condition in Lemma 1 and \( R_1 \) represented by a TEFSM and \( D_{i,j}, \forall i, j, 1 \leq i, j \leq n \). Note that \( D_{i,j} = 0, \forall i, 1 \leq i \leq n \).

- **Output:** Restricted service specification \( SS^* \) in a TEFSM

For each state \( s \) with \( |IN(s)| > 0 \) in \( SS \) do the following:

1. Let \( IN(s) = \{\{u_1, e^{in_1}, s\}, \ldots, \{u_k, e^{in_k}, s\}\} \) and \( OUT(s) = \{(s, f^{out_1}, v_1)\}, \ldots, (s, f^{out_n}, v_n)\} \).

2. For each \( i, 1 \leq i \leq n \), where \( n \) is the number of the protocol entities, do the following:
\[ \text{Lemma 3} \quad \text{Let } PS_i, 1 \leq i \leq n, \text{ be the derived protocol specification from } SS \text{ and } SS^* \text{ be the restricted service specification of } SS. \text{ Then every valid sequence } \sigma \text{ in } \{SS\} \text{ is a merged sequence from } \{SS_1, \ldots, SS_n\}, \text{ where } SS_1 \text{ preserves the order of events as specified in } PS_i, 1 \leq i \leq n. \text{ Moreover, } \{SS^*_1, 1 \leq i \leq n\}, \text{ and } \{SS^*\} \text{ is maximal in the sense that any extension of a time interval in } SS^* \text{ might be able to generate sequences which are not in } \{PS_i, 1 \leq i \leq n\} \text{ under some specific delay constraints.} \]

**Proof:** Since SS and SS* are equivalent if timing is ignored, we know that it suffices to show that \{SS^*\} \subseteq \{PS_i, 1 \leq i \leq n\} to guarantee that the order of events specified in SS is preserved in each PS_i, 1 \leq i \leq n. We first show that \{SS^*\} \subseteq \{PS_i, 1 \leq i \leq n\}. The proof is by induction on \(|\sigma|\), where \(\sigma\) is a valid sequence in SS*. Base Case \(|\sigma| = 1\). It is easy to see that \(\sigma\) is a merged sequence from \{PS_i\}, where \(\sigma = \langle a', t \rangle \). Induction Hypothesis (IH for short) Assume the claim holds for \(|\sigma| = k \geq 0\). Let \(\sigma' = \sigma < a', t >, |\sigma| = k\), be a valid sequence in SS*. By IH, we know that for each \(j, 1 \leq j \leq n\), \(\sigma_j\) is a valid sequence in \(PS_j\). We show that \(\sigma |_i < a', t >\) is also a valid sequence in \(PS_i\), where we assume for the sake of the proof that \(PS_i\) is the protocol specification obtained in Step 3, i.e., one with \(\epsilon\) transitions. Note that \(\sigma |_j = \sigma' |_j\), for \(j \neq i, 1 \leq j \leq n\). Let \(s^i = \text{head}(a')\). Since \(\sigma'\) is a valid sequence in SS*, it is clear that \(a'\) had been executable at \(\text{last}(\sigma)\). We let \(\text{last}(\sigma |_i) = \text{first}(\sigma), \text{ if } |\sigma |_i = \epsilon\). By Step 2 of the algorithm Synthesis, the transitions after \(\text{last}(\sigma |_i)\) through \(\text{last}(\sigma)\), if any, in \(PS_i\) would be \(\epsilon\) transitions. It is not hard to verify that, by investigating Step 1 and 2 of the algorithm Synthesis, the sequence \(\sigma\) would be able to lead \(PS_i\) into the state \(s\) and moreover \(s\) is reachable from either \(\text{head}(\text{last}(\sigma |_i))\), if any, or the start state of \(PS_i\), otherwise, via only \(\epsilon\) send, and/or receive transitions. Next we show that the inequalities \(t(s) + min_{oi}^{PS_i} + d_{inj,i} \leq t \leq t(s) + max_{oi}^{PS_i} + d_{inj,i}\) hold regardless of the value of the actual delay \(d_{inj,i}\) as long as \(0 < d_{inj,i} \leq D_{inj,i}\), where \([min_{oi}^{PS_i}, max_{oi}^{PS_i}]\) is the time interval associated with \(a'\) in \(PS_i\) and \(< e^{inj}, t_e >\) is an incoming event of \(s\) such that \(t(s) = t_e\). Note that \([min_{a'}^{SS^*}, max_{a'}^{SS^*}] = [min_{a'}, max_{a'} - max_{1 \leq j \leq k}(D_{inj,i})]\). Since \(\sigma'\) is a valid sequence in SS*, we have that \(t(s) + min_{a'}^{SS^*} \leq t \leq t(s) + max_{a'}^{SS^*}\), where \([min_{a'}^{SS^*}, max_{a'}^{SS^*}] = \bigcap_{0 < d_{\text{inj},i} \leq D_{\text{inj},i}} \{[min_{a'} + d, max_{a'} - max_{1 \leq j \leq k}(D_{inj,i}) + d]\}\) is the time interval associated with \(a'\) in SS*. We have that \(t(s) + min_{a'}^{PS_i} + d_{inj,i} = t(s) + min_{a'} + d_{inj,i} \leq t(s) + min_{a'} + D_{inj,i} \leq t(s) + max_{a'} + max_{1 \leq j \leq k}(D_{inj,i}) \leq t(s) + max_{a'}^{SS^*} \leq t.\) Similarly, \(t \leq t(s) + max_{a'}^{SS^*} < t(s) + max_{a'} - max_{1 \leq j \leq k}(D_{inj,i}) + \eta\), where the last inequality holds for any positive constant \(\eta\). Hence, by choosing \(\eta\) sufficiently small, we have \(t(s) + max_{a'}^{PS_i} + \eta \leq t(s) + max_{a'}^{PS_i} + d_{inj,i}\), which completes the other half. Therefore, since \(\sigma |_i\) is a valid sequence in \(PS_i\) (by IH), \(\sigma' |_i = \sigma |_i, \sigma' |_i < a', t >\) is also a valid sequence in \(PS_i\). To prove the maximality of \{SS^*\}, consider a sequence \(\psi\) in \{SS'\} - {SS^*}, where \(SS'\) is \(SS^*\) with a time interval in \(SS^*\) extended. We know that \(|\psi| > 1\), since any sequence in \{SS'\} with length 1 must have an unadjusted time interval, which implies that any extension of the interval would generate a sequence not in \{SS\}, a contradiction. We know that there exists a pair of events \(< e', t_e >, < f', t_f >\) in \(\psi\) such that \(e' \in IN(s), f^j \in OUT(s), i \neq j,\) and \(t(s) = t_e\) for some state \(s\) in \(SS\) and the time interval...
associated with $f^i$ is extended in $SS'$, since otherwise $\psi$ would not be in $\{SS'\} - \{SS^*\}$. Note that $e^i$ and $f^i$ might not be adjacent in $\psi$. By the algorithm Restriction, the time interval associated with the event $f^i$ in $SS^*$ is adjusted into $[l, u] = \{min_j + d, max_f - max_{x_j} \leq x \leq k \{D_{in,x,i}\} + d\}$. Assume the extended interval associated with the event $f^i$ in $SS'$ is $[l - \epsilon_1, u]$ or $[l, u + \epsilon_2]$, where $\epsilon_1$ and $\epsilon_2$ are positive constants. Then $l - \epsilon_1 = min_f + max_{x_j} \leq x \leq k \{D_{in,x,i}\} - \epsilon_1 < min_f + D_{ij}$, if $max_{x_j} \leq x \leq k \{D_{in,x,i}\} = D_{ij}$.

Hence, if the actual delay from $PS_i$ to $PS_j$ is $D_{ij}$ and $max_{x_j} \leq x \leq k \{D_{in,x,i}\} = D_{ij}$, then $< e^i, t_e >, < f^i, t_e + l - \epsilon_1 >$ would not be possible in $\{PS_i, 1 \leq i \leq n\}$. Similarly, $u + \epsilon_2 > max_{f_k} - max_{x_j} \leq x \leq k \{D_{in,x,i}\} + \epsilon_2$. Thus, if the actual delay from $PS_i$ to $PS_j$ is less than $\epsilon_2$, $< e^i, t_e >, < f^i, t_e + u + \epsilon_2 >$ would not be possible in $\{PS_i, 1 \leq i \leq n\}$, either.

By lemmas 2 and 3, we have the following theorem which proves the correctness of the algorithm Synthesis.

**Theorem 1** A derived protocol specification $PS_i, 1 \leq i \leq n$, is correct with respect to the service specification $SS$.

### 5 An Example

To demonstrate the synthesis method, we show the protocol specification after each step of the algorithm Synthesis when the service specification $SS$ in Figure 5 is given. Figure 6 (a), (b), and (c) describe the protocol specification $PS_1$ after each step of the algorithm Synthesis. After removing $\epsilon$ transitions, we have the final protocol specification $PS_1$, which is given in Figure 7 along with the final protocol specifications $PS_2$ and $PS_3$.

![Figure 5: A Service Specification SS](image)

### 6 Conclusion

We proposed a model based on EFSM that can represent concurrency, synchronization, and timing requirements explicitly, and presented a method to synthesize protocol specifications from timed service specifications based on the model. The proposed method appropriately inserts send and/or receive transitions between the events in the service specification so that the event orderings in the service specification are preserved. The time intervals associated with transitions are also adjusted by the method to incorporate the delay between protocol
Figure 6: \(PS_1\) after each step of the algorithm Synthesis

Figure 7: Protocol Specifications \(PS_1\), \(PS_2\), and \(PS_3\)

\[\text{empty}\] becomes true only if all the previous send or receive transitions are executed.
entities for synchronization. We proved that the derived protocol specification is optimal in the 

sense that any superset of the protocol specification would necessarily include specifications 

which are not attainable from the service specification under some specific delay constraints. 

We also presented a method to derive a sub specification from a service specification and a 

maximum communication delay of each channel such that the sub specification, but no superset 

of it, can be simulated by the derived protocol specification.

A formalization of logical errors in a TEFSM would be required to further investigate the 

relationship between the derived protocol specification and the service specification, as far as 

insuring the absence of design errors.

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